

ECS455: Chapter 5

OFDM

5.4 Cyclic Prefix (CP)



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Tuesday 9:30-10:30

Tuesday 13:30-14:30

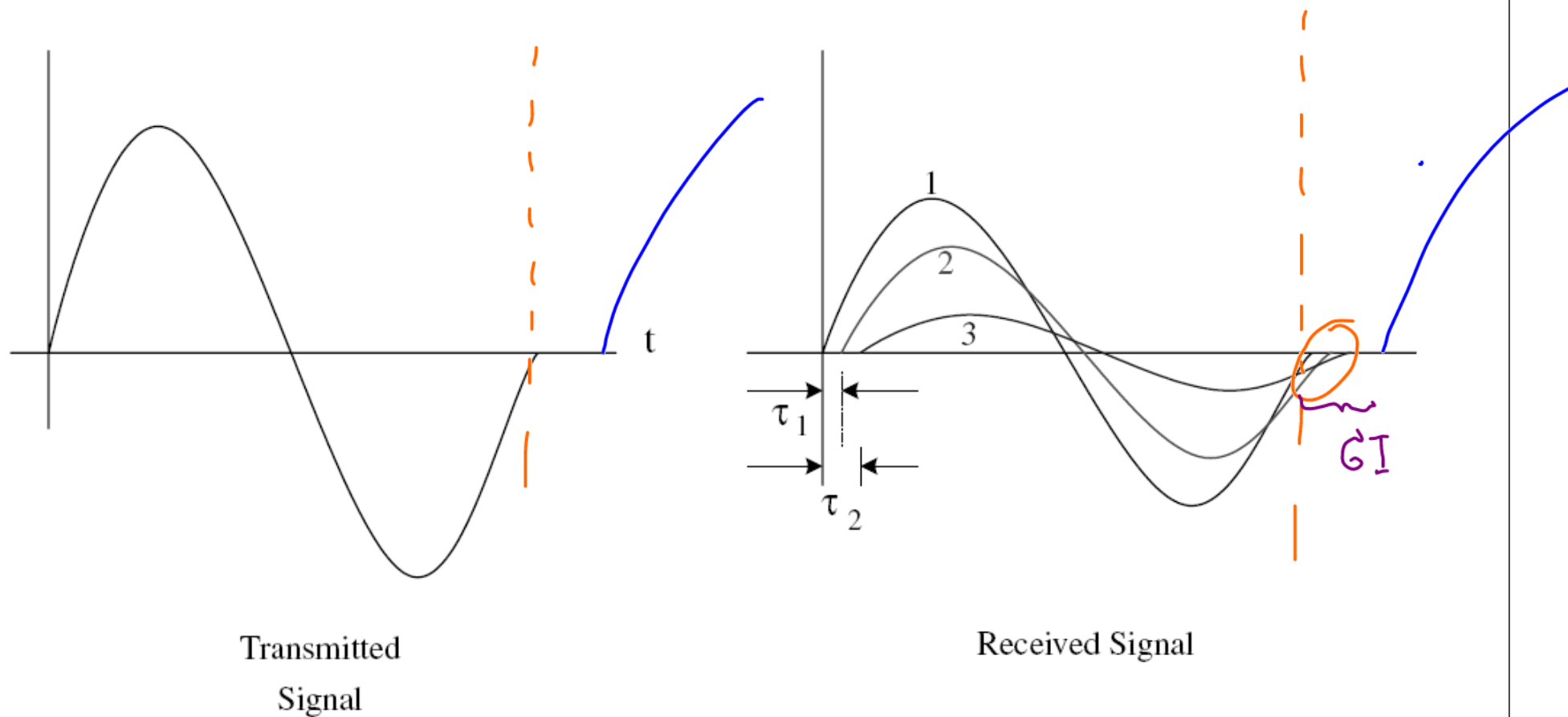
Thursday 13:30-14:30

Three steps towards modern OFDM

1. Mitigate Multipath (ISI): Decrease the rate of the original data stream via multicarrier modulation (FDM)
2. Gain Spectral Efficiency: Utilize orthogonality
3. Achieve Efficient Implementation: FFT and IFFT
 - Extra step: Completely eliminate ISI and ICI
 - Cyclic prefix

Cyclic Prefix: Motivation (1)

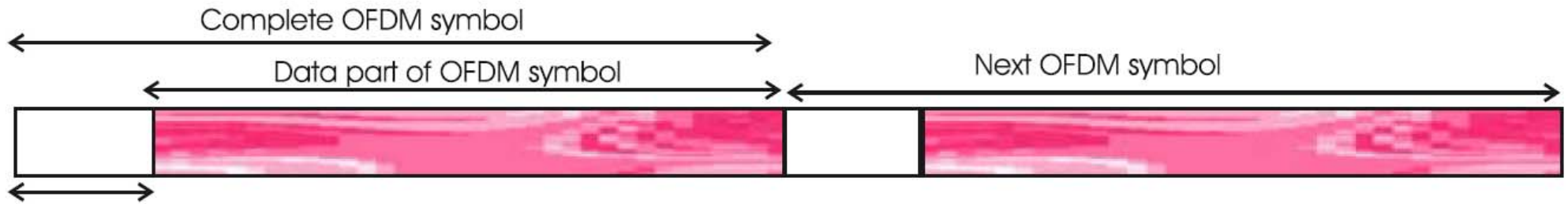
- Recall: Multipath Fading and Delay Spread



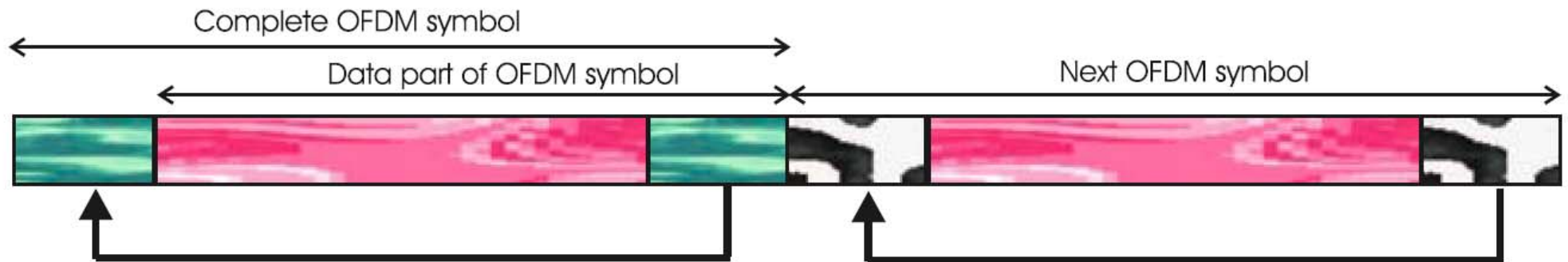
Cyclic Prefix: Motivation (2)

- OFDM uses large symbol duration T_s
 - compared to the duration of the impulse response τ_{\max} of the channel
 - to reduce the amount of ISI
- **Q:** Can we “eliminate” the multipath (**ISI**) problem?
- **A:** To reduce the ISI, add **guard interval** larger than that of the estimated delay spread.
- If the guard interval is left empty, the orthogonality of the sub-carriers no longer holds, i.e., **ICI** (inter-channel interference) still exists.
- **Solution:** To prevent **both** the **ISI** as well as the **ICI**, OFDM symbol is **cyclically extended** into the guard interval.

Cyclic Prefix



Guard Interval, $T_{CP} > \tau_{max}$
Using empty spaces as guard interval at the beginning of each symbol



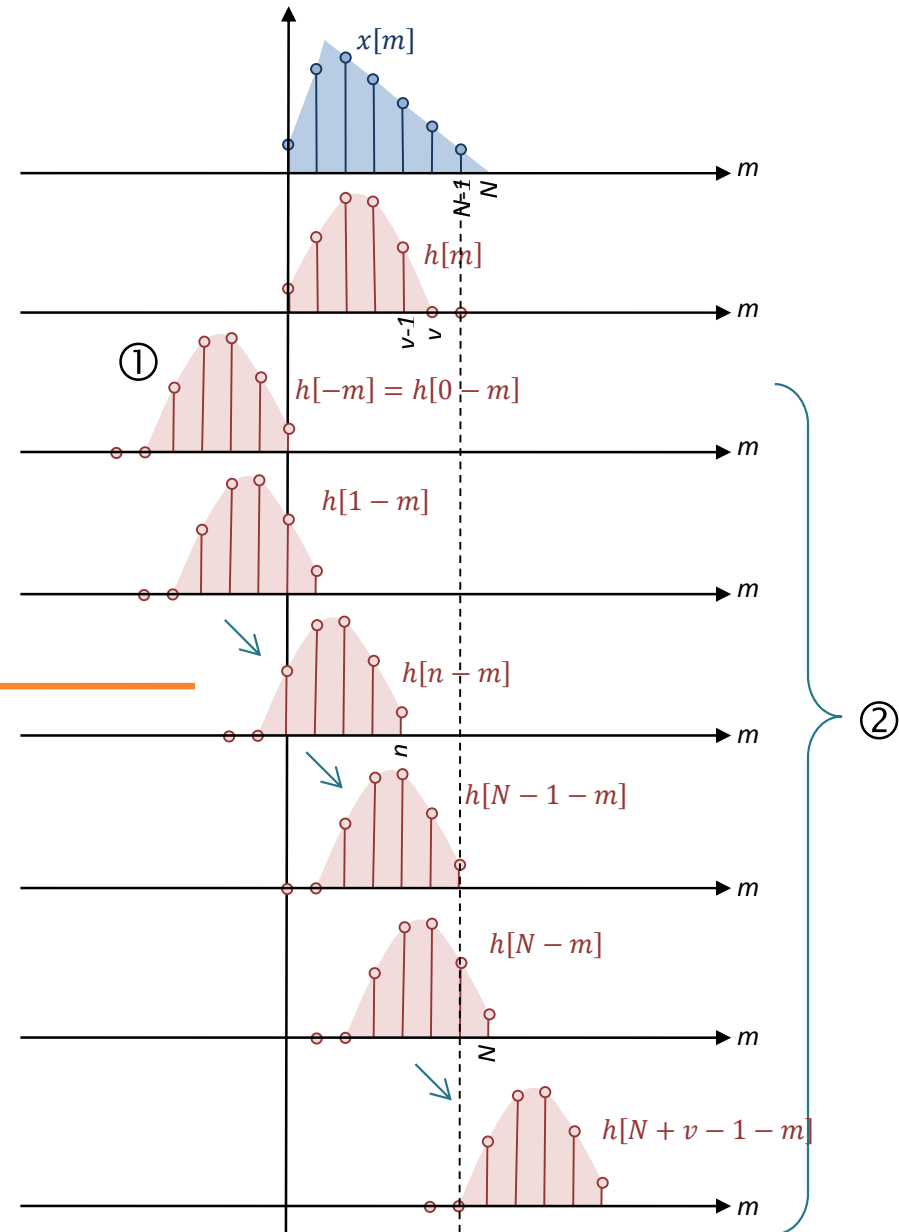
End of symbol is prepended to beginning
Guard interval still equals to T_{CP}

Using cyclic prefix:
OFDM symbol length: $T_{sym} + T_{CP}$
Efficiency: $T_{sym} / (T_{sym} + T_{CP})$

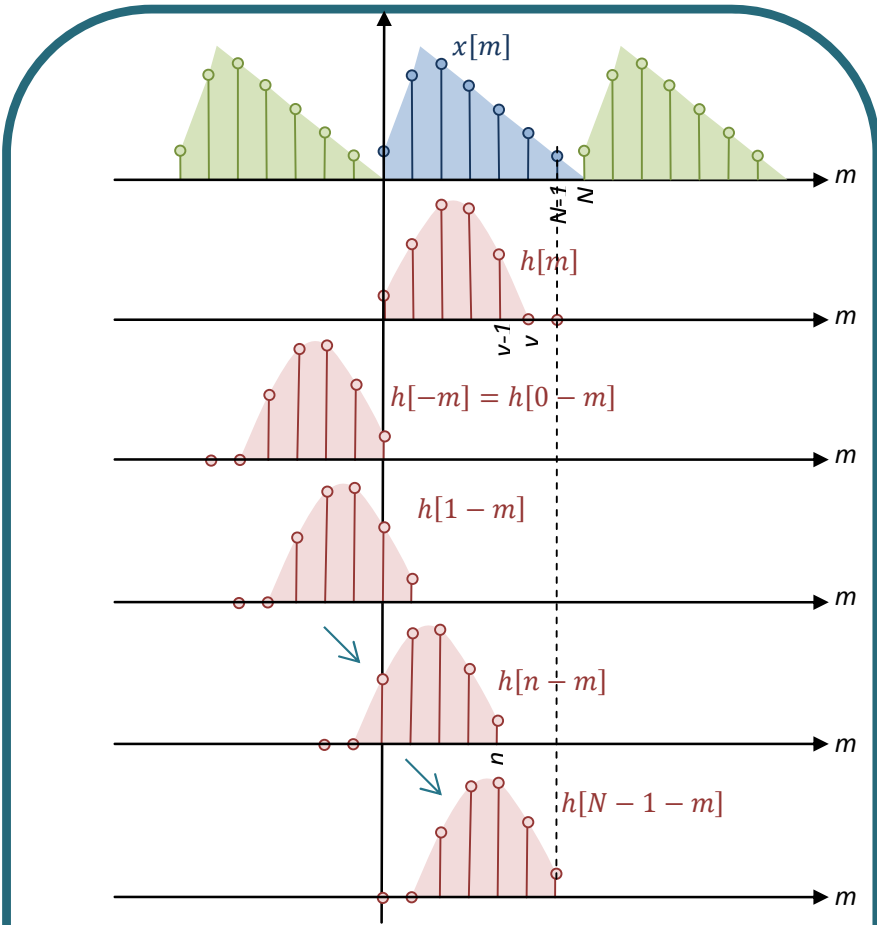
Recall: Convolution

- ① Flip
- ② Shift
- ③ Multiply (pointwise)
- ④ Add

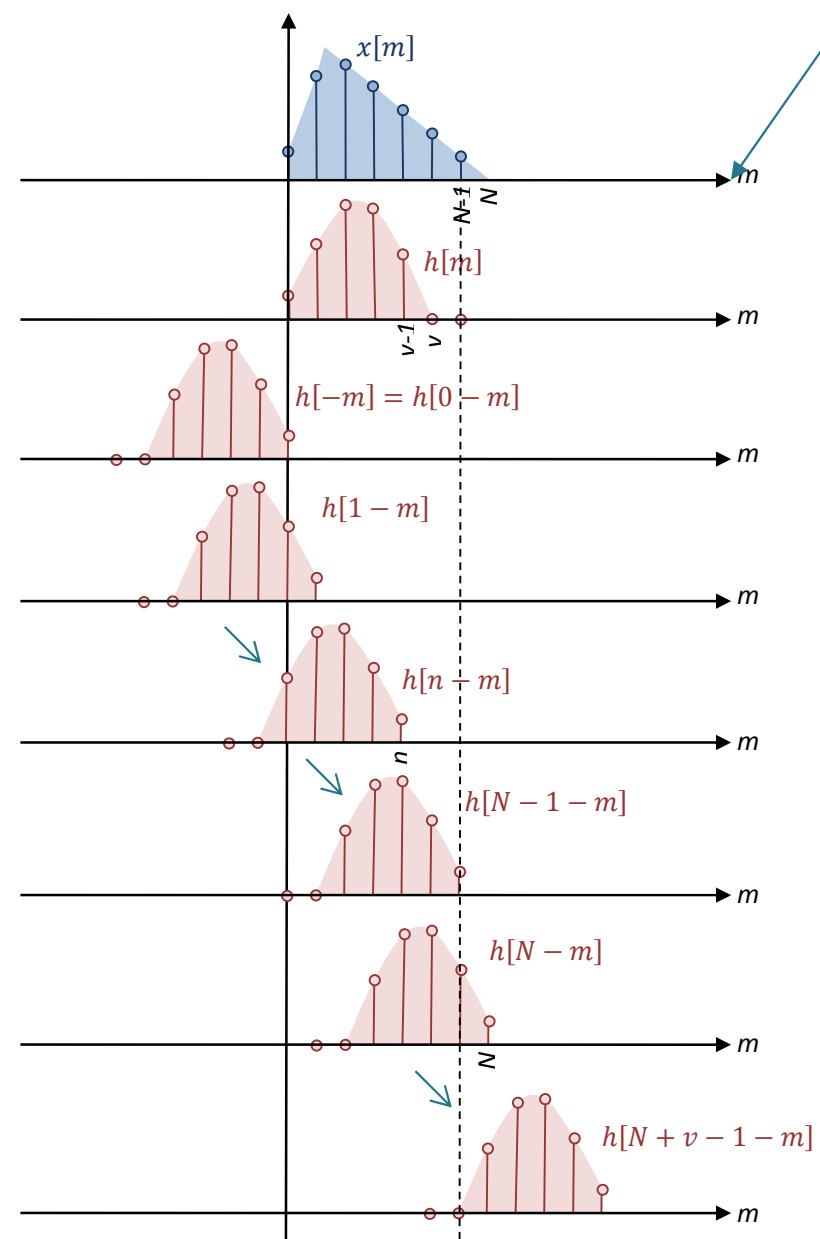
$$\{x * h\}[n] = \sum_m x[m] h[n - m]$$



Circular Convolution



(Regular Convolution)



0 Replicate x (now it looks periodic)
 5 Then, perform the usual convolution
 only on $n = 0$ to $N-1$

Circular Convolution: Examples 1

Find

$$[1 \ 2 \ 3] * [4 \ 5 \ 6]$$

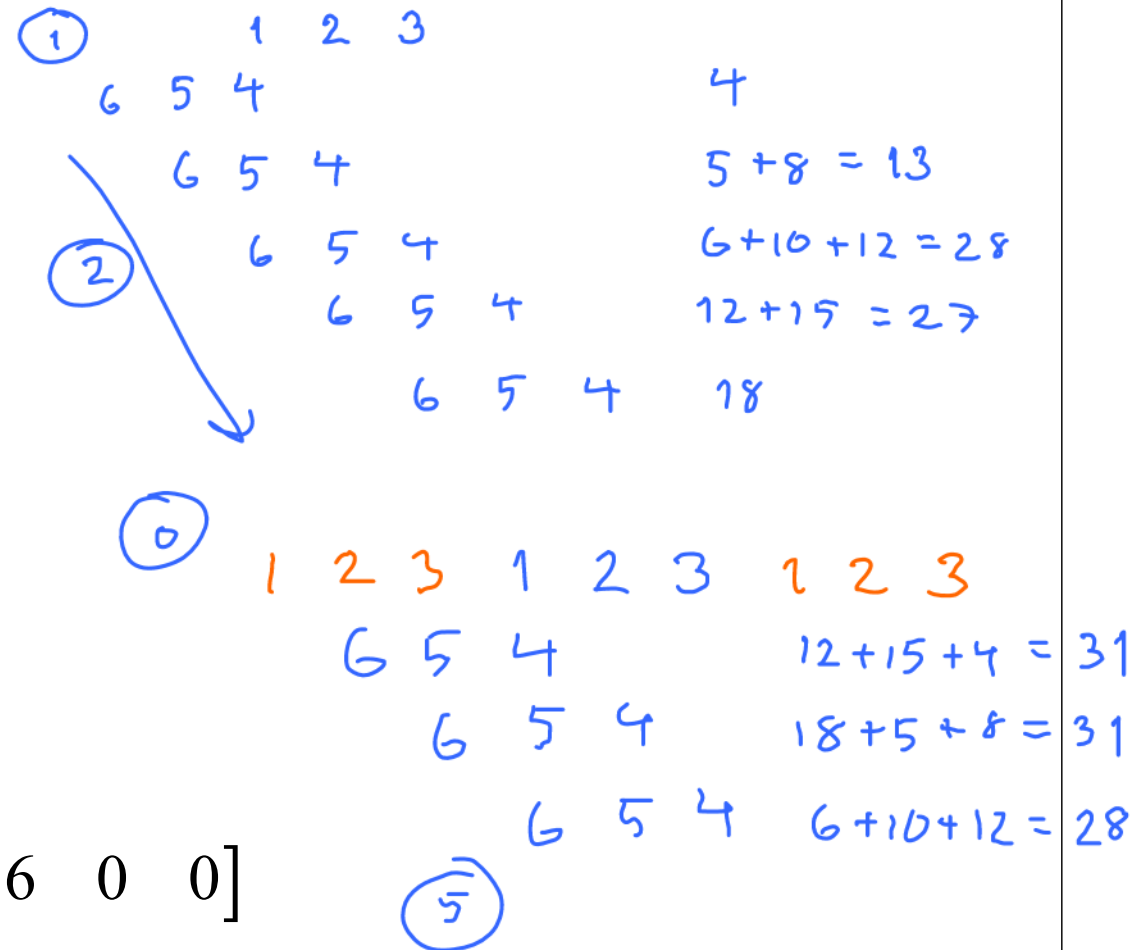
```
>> conv([1,2,3],[4,5,6])
ans =
     4     13     28     27     18
```

$$[1 \ 2 \ 3] \otimes [4 \ 5 \ 6]$$

```
>> cconv([1,2,3],[4,5,6],3)
ans =
    31     31     28
```

$$[1 \ 2 \ 3 \ 0 \ 0] \otimes [4 \ 5 \ 6 \ 0 \ 0]$$

```
>> cconv([1,2,3,0,0],[4,5,6,0,0],5)
ans =
  4.0000  13.0000  28.0000  27.0000  18.0000
```



Discussion

- *Regular convolution* of an N_1 -point vector and an N_2 -point vector gives (N_1+N_2-1) -point vector.
- *Circular convolution* is performed between two equal-length vectors. The results also has the same length.
- Circular convolution can be used to find regular convolution by **zero-padding**.
 - Zero-pad the vectors so that their length is N_1+N_2-1 .
 - Example:
$$[1 \ 2 \ 3 \ 0 \ 0] \circledast [4 \ 5 \ 6 \ 0 \ 0] = [1 \ 2 \ 3] * [4 \ 5 \ 6]$$
- In modern OFDM, we want to perform circular convolution via regular convolution.

Circular Convolution in Communication

- We want the receiver to obtain the circular convolution of the signal (channel input) and the channel.
- Q: Why?
- A:
 - **CTFT: convolution in time domain** corresponds to **multiplication in frequency domain.**
 - This fact does not hold for DFT.
 - **DFT: circular convolution** in (discrete) time domain corresponds to **multiplication** in (discrete) frequency domain.
 - We want to have multiplication in frequency domain.
 - So, we want circular convolution and not the regular convolution.
- Problem: Real channel does regular convolution.
- Solution: With **cyclic prefix**, regular convolution can be used to create circular convolution.

Example 2

$$[1 \ -2 \ 3 \ 1 \ 2] \circledast [3 \ 2 \ 1 \ 0 \ 0] = ?$$

Let's look closer at how we carry out the circular convolution operation. Recall that we replicate the x and then perform the regular convolution (for N points)

Solution:

$$\begin{array}{cccccccccccccccc}
 1 & -2 & 3 & 1 & 2 & 1 & -2 & 3 & 1 & 2 & 1 & -2 & 3 & 1 & 2 \\
 0 & 0 & 1 & 2 & 3 & & & & & & & & & & \\
 0 & 0 & 1 & 2 & 3 & & & & & & & & & & \\
 & 0 & 0 & 1 & 2 & 3 & & & & & & & & & \\
 & & 0 & 0 & 1 & 2 & 3 & & & & & & & & \\
 & & & 0 & 0 & 1 & 2 & 3 & & & & & & & \\
 & & & & 0 & 0 & 1 & 2 & 3 & & & & & &
 \end{array}$$

$$1 \times 1 + 2 \times 2 + 1 \times 3 = 1 + 4 + 3 = 8$$

$$2 \times 1 + 1 \times 2 + (-2) \times 3 = 2 + 2 - 6 = -2$$

$$1 \times 1 + (-2) \times 2 + 3 \times 3 = 1 - 4 + 9 = 6$$

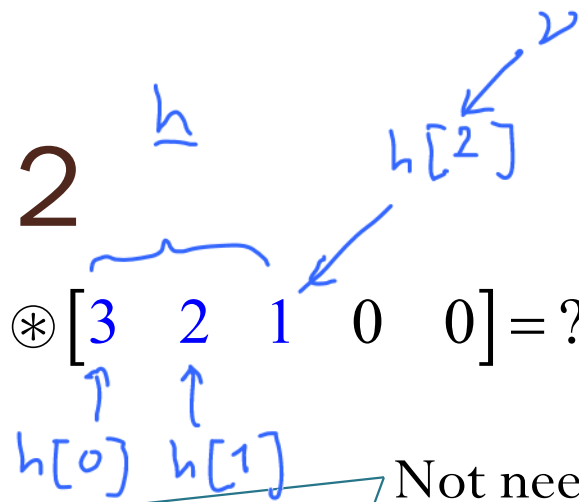
$$(-2) \times 1 + 3 \times 2 + 1 \times 3 = -2 + 6 + 3 = 7$$

$$3 \times 1 + 1 \times 2 + 2 \times 3 = 3 + 2 + 6 = 11$$

$$[1 \ -2 \ 3 \ 1 \ 2] \circledast [3 \ 2 \ 1 \ 0 \ 0] = [8 \ -2 \ 6 \ 7 \ 11]$$

Example 2

$$[1 \ -2 \ 3 \ 1 \ 2] \circledast [3 \ 2 \ 1 \ 0 \ 0] = ?$$



Observation: We don't need to replicate the x indefinitely. Furthermore, when h is shorter than x , we need only a part of one replica.

1 -2 3

1 2 1 -2 3 1 2

1 -2 3 1 2

0 0 1 2 3

$$1 \times 1 + 2 \times 2 + 1 \times 3 = 1 + 4 + 3 = 8$$

0 0 1 2 3

$$2 \times 1 + 1 \times 2 + (-2) \times 3 = 2 + 2 - 6 = -2$$

0 0 1 2 3

$$1 \times 1 + (-2) \times 2 + 3 \times 3 = 1 - 4 + 9 = 6$$

0 0 1 2 3

$$(-2) \times 1 + 3 \times 2 + 1 \times 3 = -2 + 6 + 3 = 7$$

0 0 1 2 3

$$3 \times 1 + 1 \times 2 + 2 \times 3 = 3 + 2 + 6 = 11$$

$$[1 \ -2 \ 3 \ 1 \ 2] \circledast [3 \ 2 \ 1 \ 0 \ 0] = [8 \ -2 \ 6 \ 7 \ 11]$$

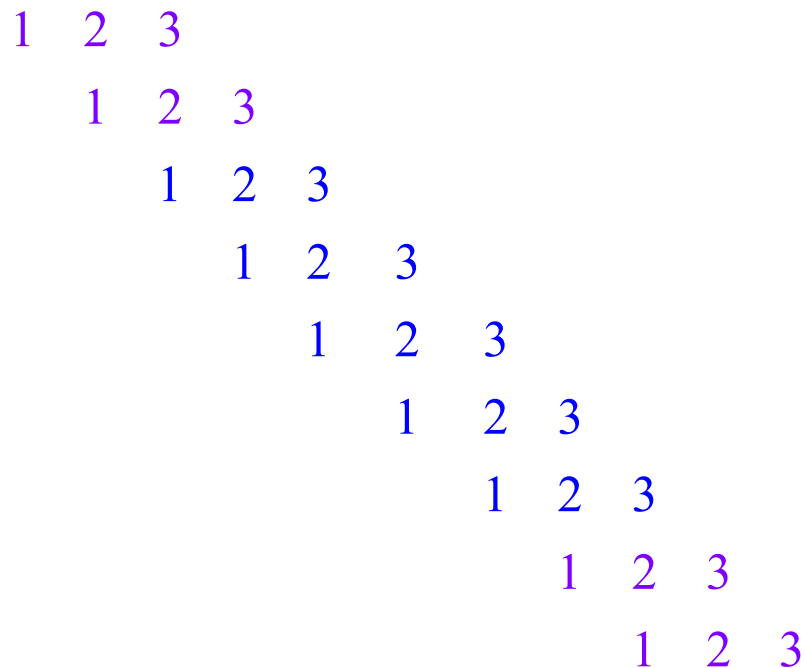
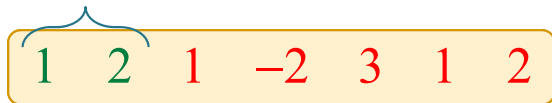
Example 2

Try this: use only the necessary part of the replica and then convolute (regular convolution) with the channel.

$$[1 \ 2 \ 1 \ -2 \ 3 \ 1 \ 2] * [3 \ 2 \ 1] = ?$$

Copy the last v samples of the symbols at the **beginning** of the symbol.

This partial replica is called the **cyclic prefix**.



$$\begin{aligned}
 & 1 \times 3 = 3 \\
 & 1 \times 2 + 2 \times 3 = 2 + 6 = 8 \\
 & 1 \times 1 + 2 \times 2 + 1 \times 3 = 1 + 4 + 3 = 8 \\
 & 2 \times 1 + 1 \times 2 + (-2) \times 3 = 2 + 2 - 6 = -2 \\
 & 1 \times 1 + (-2) \times 2 + 3 \times 3 = 1 - 4 + 9 = 6 \\
 & (-2) \times 1 + 3 \times 2 + 1 \times 3 = -2 + 6 + 3 = 7 \\
 & 3 \times 1 + 1 \times 2 + 2 \times 3 = 3 + 2 + 6 = 11 \\
 & 1 \times 1 + 2 \times 2 = 1 + 4 = 5 \\
 & 2 \times 1 = 2
 \end{aligned}$$

Junk!

Example 2

- We now know that

$$\underbrace{[1 \ 2 \ 1 \ -2 \ 3 \ 1 \ 2]}_{\text{Cyclic Prefix}} * [3 \ 2 \ 1] = [3 \ 8 \ 8 \ -2 \ 6 \ 7 \ 11 \ 5 \ 2]$$

$$[1 \ -2 \ 3 \ 1 \ 2] \circledast [3 \ 2 \ 1 \ 0 \ 0]$$

- Similarly, you may check that

$$\underbrace{[-2 \ 1 \ 2 \ 1 \ -3 \ -2 \ 1]}_{\text{Cyclic Prefix}} * [3 \ 2 \ 1] = [-6 \ -1 \ 6 \ 8 \ -5 \ -11 \ -4 \ 0 \ 1]$$

$$[2 \ 1 \ -3 \ -2 \ 1] \circledast [3 \ 2 \ 1 \ 0 \ 0]$$

Example 3

- We know, from Example 2, that

$$[\text{1 2 1 -2 3 1 2}] * [\text{3 2 1}] = [\text{3 8 8 -2 6 7 11 5 2}]$$

And that

$$[\text{-2 1 2 1 -3 -2 1}] * [\text{3 2 1}] = [\text{-6 -1 6 8 -5 -11 -4 0 1}]$$

- Check that

$$\begin{aligned} & [\text{1 2 1 -2 3 1 2 0 0 0 0 0 0 0}] * [\text{3 2 1}] \\ = & [\text{3 8 8 -2 6 7 11 5 2 0 0 0 0 0}] \end{aligned}$$

and

$$\begin{aligned} & [\text{0 0 0 0 0 0 0 0 -2 1 2 1 -3 -2 1}] * [\text{3 2 1}] \\ = & [\text{0 0 0 0 0 0 0 0 -6 -1 6 8 -5 -11 -4 0 1}] \end{aligned}$$

Example 4

- We know that

$$\begin{bmatrix} 1 & 2 & 1 & -2 & 3 & 1 & 2 \end{bmatrix} * \begin{bmatrix} 3 & 2 & 1 \end{bmatrix} = \begin{bmatrix} 3 & 8 & 8 & -2 & 6 & 7 & 11 & 5 & 2 \end{bmatrix}$$

$$\begin{bmatrix} -2 & 1 & 2 & 1 & -3 & -2 & 1 \end{bmatrix} * \begin{bmatrix} 3 & 2 & 1 \end{bmatrix} = \begin{bmatrix} -6 & -1 & 6 & 8 & -5 & -11 & -4 & 0 & 1 \end{bmatrix}$$

- Using Example 3, we have

$$\begin{bmatrix} 1 & 2 & 1 & -2 & 3 & 1 & 2 & -2 & 1 & 2 & 1 & -3 & -2 & 1 \end{bmatrix} * \begin{bmatrix} 3 & 2 & 1 \end{bmatrix}$$

$$= \left(\begin{array}{l} \begin{bmatrix} 1 & 2 & 1 & -2 & 3 & 1 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \\ + \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & -2 & 1 & 2 & 1 & -3 & -2 & 1 \end{bmatrix} \end{array} \right) * \begin{bmatrix} 3 & 2 & 1 \end{bmatrix}$$

$$= \begin{array}{l} \begin{bmatrix} 3 & 8 & 8 & -2 & 6 & 7 & 11 & 5 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \\ + \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & -6 & -1 & 6 & 8 & -5 & -11 & -4 & 0 & 1 \end{bmatrix} \end{array}$$

$$= \begin{bmatrix} 3 & 8 & 8 & -2 & 6 & 7 & 11 & -1 & 1 & 6 & 8 & -5 & -11 & -4 & 0 & 1 \end{bmatrix}$$

Putting results together...

- Suppose $\underline{x}^{(1)} = [1 \ -2 \ 3 \ 1 \ 2]$ and $\underline{x}^{(2)} = [2 \ 1 \ -3 \ -2 \ 1]$
- Suppose $\underline{h} = [3 \ 2 \ 1]$
- At the receiver, we want to get
 - $[1 \ -2 \ 3 \ 1 \ 2] \otimes [3 \ 2 \ 1 \ 0 \ 0] = [8 \ -2 \ 6 \ 7 \ 11]$
 - $[2 \ 1 \ -3 \ -2 \ 1] \otimes [3 \ 2 \ 1 \ 0 \ 0] = [6 \ 8 \ -5 \ -11 \ -4]$
- We transmit $[\underbrace{1 \ 2}_{\text{Cyclic prefix}} \ 1 \ -2 \ 3 \ 1 \ 2 \ \underbrace{-2 \ 1}_{\text{Cyclic prefix}} \ 2 \ 1 \ -3 \ -2 \ 1]$.

- At the receiver, we get

$$[\ 1 \ 2 \ 1 \ -2 \ 3 \ 1 \ 2 \ -2 \ 1 \ 2 \ 1 \ -3 \ -2 \ 1] * [3 \ 2 \ 1]$$

$$= [\ 3 \ 8 \ 8 \ -2 \ 6 \ 7 \ 11 \ -1 \ 1 \ 6 \ 8 \ -5 \ -11 \ -4 \ 0 \ 1]$$

Junk! To be thrown away by the receiver.

Circular Convolution: Key Properties

- Consider an N -point signal $x[n]$
- **Cyclic Prefix (CP) insertion:** If $x[n]$ is extended by copying the last ν samples of the symbols at the beginning of the symbol:

$$\hat{x}[n] = \begin{cases} x[n], & 0 \leq n \leq N-1 \\ x[n+N], & -\nu \leq n \leq -1 \end{cases}$$

- Key Property 1:

$$\{h \circledast x\}[n] = (h * \hat{x})[n] \text{ for } 0 \leq n \leq N-1$$

- Key Property 2:

$$\{h \circledast x\}[n] \xrightarrow{\text{FFT}} H_k X_k$$

$$\mathbf{h} = (h[0], h[1], h[2], \dots, h[\nu])$$

$$\mathbf{H} = \text{FFT}(\tilde{\mathbf{h}})$$

↑ zero-padded to length N

OFDM with CP for Channel w/ Memory

- To send N samples $\mathbf{S} = (S_0, S_1, \dots, S_{N-1})$
- First apply IFFT with scaling by \sqrt{N} : $\tilde{\mathbf{s}} = \sqrt{N} \text{IFFT}(\mathbf{S})$
- Then, add **cyclic prefix**

$$\mathbf{x} = [\tilde{s}[N - \nu], \dots, \tilde{s}[N - 1], \tilde{s}[0], \dots, \tilde{s}[N - 1]]$$

- This is inputted to the channel
- The channel output is $\mathbf{y} = \mathbf{x} * \mathbf{h}$ which can be viewed as

$$\mathbf{y} = [p[N - \nu], \dots, p[N - 1], r[0], \dots, r[N - 1]]$$

- Remove cyclic prefix to get \mathbf{r} . (We know that $\mathbf{r} = \tilde{\mathbf{s}} \circledast \mathbf{h}$.)
- Then apply FFT with scaling by $1/\sqrt{N}$: $\tilde{\mathbf{R}} = \frac{1}{\sqrt{N}} \text{FFT}(\mathbf{r})$
- By circular convolution property of DFT,

$$\mathbf{r} = \tilde{\mathbf{s}} \circledast \mathbf{h} \longrightarrow R_k = H_k \tilde{S}_k \longrightarrow \tilde{R}_k = H_k S_k \longrightarrow S_k = \frac{\tilde{R}_k}{H_k}$$

No ICI!

MATLAB Example (1/2)

```
S = [1 -1 2 4 5 -1 2 -3]; % data stream  
h = [1 0.3 0.1];
```

```
% OFDM transmitter
```

```
N = 4;
```

```
n = length(S)/N;
```

```
St = (reshape(S,N,n)).';
```

```
st = (sqrt(N))*ifft(St,[],2);
```

```
% Number of data symbols per OFDM symbol
```

```
% Number of data blocks
```

```
% Reshape stream to matrix for
```

```
% easier addition of cyclic prefix
```

```
% Calculate the IFFT with scaling
```

```
St =  
    1    -1     2     4  
    5    -1     2    -3
```

IFFT
w/ scaling
(row-wise)

```
st =  
 3.0 + 0.0i  -0.5 - 2.5i  0.0 + 0.0i  -0.5 + 2.5i  
 1.5 + 0.0i  1.5 + 1.0i  5.5 + 0.0i  1.5 - 1.0i
```

```
v = length(h)-1;
```

```
xt = [st(:,(N-(v-1)):N) st];
```

```
% Add Cyclic Prefix
```

```
x = (reshape(xt.',(N+v)*n,1)).'; % Reshape back to stream
```

```
x =  
 0.0 + 0.0i  -0.5 + 2.5i  3.0 + 0.0i  -0.5 - 2.5i  0.0 + 0.0i  -0.5 + 2.5i  5.5 + 0.0i  1.5 - 1.0i  1.5 + 0.0i  
 1.5 + 1.0i  5.5 + 0.0i  1.5 - 1.0i
```

MATLAB Example (2/2)

```

%-----
% Convolve with channel
y = conv(x,h);
H = fft([h zeros(1,N-v-1)]);
%-----

```

```

y =
  0.00 + 0.00i  -0.50 + 2.50i  2.85 + 0.75i  0.35 - 2.25i  0.15 - 0.75i
 -0.55 + 2.25i  5.35 + 0.75i  3.10 - 0.75i  2.50 - 0.30i  2.10 + 0.90i
 6.10 + 0.30i  3.30 - 0.90i

```

```

% OFDM receiver
y = y(1:((N+v)*n));
yt = reshape(y,(N+v),n).';

r = yt(:,v+1:v+N);
Rt = (1/sqrt(N))*fft(r,[],2);

```

```

% Reshape matrix for easier
% removal of cyclic prefix
% Eliminate junk (cyclic prefix)
% Calculate the FFT with scaling

```

```

r =
  2.85 + 0.75i  0.35 - 2.25i  0.15 - 0.75i  -0.55 + 2.25i
  2.50 - 0.30i  2.10 + 0.90i  6.10 + 0.30i  3.30 - 0.90i

```

FFT
w/ scaling
(row-wise)

```

Rt =
  1.4 + 0.0i  -0.9 + 0.3i  1.6 + 0.0i  3.6 + 1.2i
  7.0 + 0.0i  -0.9 + 0.3i  1.6 + 0.0i  -2.7 - 0.9i

```

```

% "Equalization"
S_hatt = zeros(size(Rt));
for i=1:length(H)
    S_hatt(:,i) = Rt(:,i)/H(i);
end
S_hat = reshape(S_hatt.',1,N*n)

```

% Divide Rk (the kth column) by Hk

```

Shatt =
  1  -1  2  4
  5  -1  2  -3

```

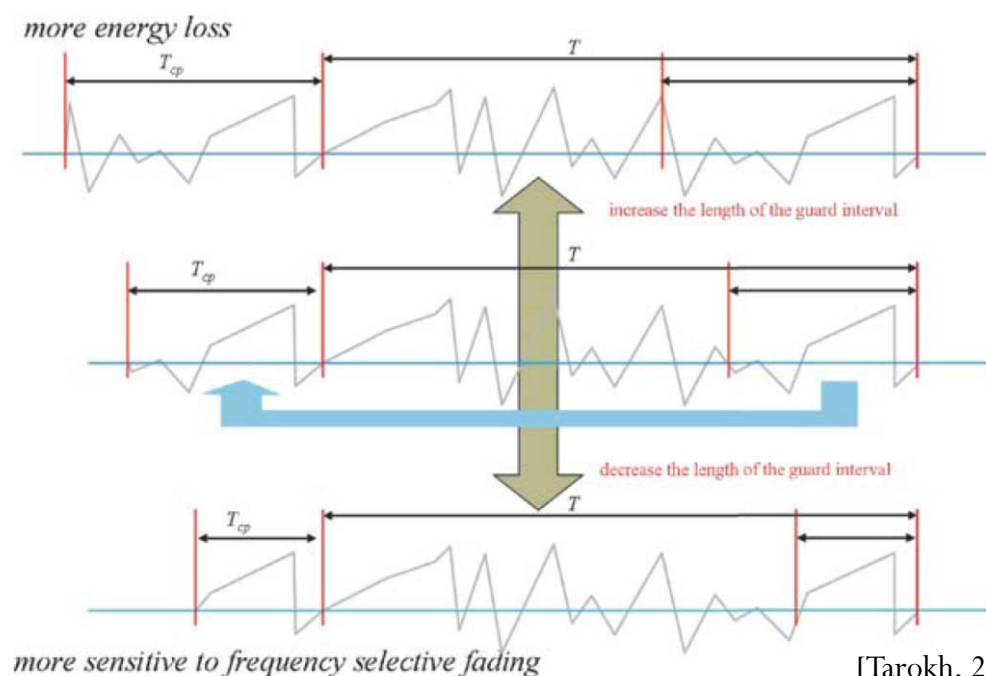
```

S_hat = [1 -1 2 4 5 -1 2 -3]

```

OFDM System Design: CP

- A good ratio between the CP interval and symbol duration should be found, so that all multipaths are resolved and not significant amount of energy is lost due to CP.
- As a thumb rule, the CP interval must be two to four times larger than the root mean square (RMS) delay spread.



Summary

- The CP at the beginning of each block has two main functions.
- As guard interval, it prevents contamination of a block by ISI from the previous block.
- It makes the received block appear to be periodic of period N .
 - Turn regular convolution into circular convolution
 - Point-wise multiplication in the frequency domain

Reference

- A. Bahai, B. R. Saltzberg, and M. Ergen, *Multi-Carrier Digital Communications: Theory and Applications of OFDM*, 2nd ed., New York: Springer Verlag, 2004.

